

Editorial for Volume 9 Number 3: Models, Resources and Activities of Project Scheduling Problems

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Abstract—The Project Scheduling Problem (PSP) is a generic name given to a whole class of problems in which the best form, time, resources and costs for project scheduling are necessary. The PSP is an application area related to the project management. This paper aims at being a guide to understand PSP by presenting a survey of the general parameters of PSP: the Resources (those elements that realize the activities of a project), and the Activities (set of operations or own tasks of a person or organization); the mathematical models of the main variants of PSP and the algorithms used to solve the variants of the PSP. The project scheduling is an important task in project management. This paper contains mathematical models, resources, activities, and algorithms of project scheduling problems. The project scheduling problem has attracted researchers of the automotive industry, steel manufacturer, medical research, pharmaceutical research, telecommunication, industry, aviation industry, development of the software, manufacturing management, innovation and technology management, construction industry, government project management, financial services, machine scheduling, transportation management, and others. The project managers need to finish a project with the minimum cost and the maximum quality.

Keywords— Project Management; Manufacturing Management, Technology Management.

1. INTRODUCTION

The Project management is considered as The planning, monitoring and control of all aspects of a project and the motivation of all those involved in it to achieve the project objectives on time and to the specified cost, quality and performance [1]. A project is a temporary and unique effort that, with a set of resources, it looks forward to satisfying specific objectives in a period with certain time [2]. The application areas of the project management usually are defined in terms of: technical elements (development of software, pharmaceutical drugs or civil engineering, Production systems planning), elements of the administration (Project Scheduling Problems, Manufacturing Management [3], Technology Management, contracts with the government or development of new products), and groups of industry (industrial engineering, automobiles, chemicals or financial services).

This paper aims at being a guide to understand the Project Scheduling Problems (PSP) by presenting a survey of the general parameters of the PSP, the algorithms used to solve the problems and the differences of the variants of the problems. Section 2 presents the definition, classification and the general parameters of the PSP, Section 3 the variants of the PSP, Section 4 the algorithms used to solve the PSP, and the last section presents the conclusions.

2. PROJECT SCHEDULING PROBLEM

Kalelkar and Critical Path Method (CPM) is a project modelling technique and it is commonly used with all forms of projects, including construction, aerospace and defense, software development, research projects, product development, engineering, and plant maintenance, among others [4]. Any project with interdependent activities can apply this method of mathematical analysis. Although the original CPM program and approach is no longer used, the term is generally applied to any approach used to analyse a project network logic diagram. The concept of CPM known (at the moment) as Project Scheduling Problem.

The Project Scheduling Problem (PSP) is a generic name given to a whole class of problems in which the best form, time, resources and costs for project scheduling are necessary. The PSP is classified according to the optimization objective by which they were created, for example, PSP of general way looks to optimize: time, project cost and resource usage.

The general parameters of PSP are:

1. The resources are those elements that realize the activities of a project. Examples of discrete resources are: machines, tools, workers; continuous resources include: energy, liquids and money. The diverse types of resources are:
 - Renewable, nonrenewable and doubly constrained resources (Slowinski, 1980; Weglarz, 1980). The renewable resources are represented by labour, machinery, equipment; the nonrenewable resources usually are represented by money; a doubly constrained resource can be either consumed (money, raw materials) or used (blades, cartridges) by activity during its execution.
 - Preemptable or non-preemptable [5,6]. The preemptable resources are those resources used for the processing of current activity, allotted to another activity, and then returned to the previously interrupted activity whose processing may be resumed; the resources without this property are called non-preemptable resources.
 - Reusable resources [7]. The reusable resources represent resources which act like renewable (recycling) but are consumed by little in each period they are used. Examples of reusable resources are tools with single or multiple cutting edges. The recycling takes time, and in consequence, each reusable resource becomes unavailable for some time periods.
 - Dedicated resources [8]. The dedicated resources can be assigned to only one activity at a time.
 - Spatial resources [9]. The spatial resources are resources used by a group of activities from the start of the first activity till the completion of the last activity of the group. Examples of spatial resources include dry docks of a shipyard, shop floor space, rooms, pallets, containers.
 - Partially (non) renewable resources [10]. The partially (non) renewable resources are resources with availabilities (subsets of periods or time intervals).

- Unary resources (Baptiste et al., 1999). The unary resources are a special case of renewable resources with availabilities limited to one unit per period.
- Complementary resources [11]. The complementary resources are used for setting up resources and not used for the processing of activities. An example of a complementary resource is a set of workers who are capable of and responsible for properly setting up other resources like specialized machines, computers, or robots.
- Cumulative resources [12]. The concept of cumulative resources is a generalization of the concept of nonrenewable resources. The cumulative resources are considered when the availability of a resource in a given time period cumulatively depends on the utilization (depletions and replenishments) of this resource during the previous time periods. Examples of such resources are investment capital, storage facilities and intermediate products.
- Multi-skill resources [13]. The multi-skill resources are able (flexibility) to be allocated for different kinds of resource requirements. Each resource usually has more than one skill, it can still be assigned to only one activity at a time. Examples are staff members (each of them has his own set of skills), multi-skill machines, tools, robots. A multi-skill resource has the functionality of several renewable resources.
- Synchronizing resources [14]. The Synchronizing resources have the ability to ensure a simultaneous start of a set of activities to which these resources are allocated. Synchronizing resources are a special category of renewable resources.
- Allocating resources [14]. The allocatable resources can be viewed as renewable resources which require some setup operations before they are ready to process a given activity. If an activity A_i requires a number of units of such a resource, they have to be allocated to this activity by another activity A_k (so-called allocating activity) starting no later than activity A_i . In other words, activity A_k prepares a given number of units of an allocatable resource for processing activity A_i . For example, a resource is equipment that has to be installed each time before it is used for executing some activity A_i . If this setup operation also requires some scarce renewable resources (like staff or tools), it can be modelled as an activity A_k .
- Adjacent resources [16,17]. The adjacent resources are resources with a physical location of a particular resource among resource (of the same type) are given, and it is important for activity processing. Examples of adjacent resources include adjacent parts of dry docks in a shipyard and check-in desks at airports, as well as other types of resources like processors in multi-processor or grid environments.
- Changeover resources [18]. The changeover resources are renewable resources with a setup (sequence-dependent) or must be changed over when passing from one activity A_i to another activity A_j . For example: a project with several geographically distributed locations and some renewable resources can be used in each location, but in order to be put into service for processing activities in a given location, it is necessary to transport units of such resources from another location. The time needed for transport is treated as a setup (changeover) time.

- Auxiliary resources [19]. This is a similar concept of complementary resources of Artigues and Roubellat (2001).
 - Heterogeneous resources [20]. The heterogeneous resources are renewable resources and multi-skills resources. An individual of a heterogeneous resource has the same skill, but some individuals may be less or more skilled than the others. The time needed to perform an activity and the quality of the performed work depend on the skill level of the individuals assigned to the activity.
2. The activities are defined as a set of operations or own tasks of a person or organization. The activities are characterized by resource requirements, activity processing model, and precedence constraints with other activities. The parameters related to the activities [21] are:
- The resource requirement (resource demand, resource request) is the amount of a resource needed to execute a given activity.
 - Activity processing model describes a processing characteristic of an activity as a function of resource amounts allotted for this activity.
 - Time Parameters. The duration of a project activity can be determined by the time parameters: a ready time, a due date or a deadline. Ready time (also known as a release date) is a time at which an activity is ready for processing. In many projects scheduling problems ready times for all activities are identical to zero. Due date specifies a time limit for completing an activity. Usually, penalty costs per unit are considered for activities completed either before/after their due dates. The deadline can be defined similarly to the due date. The main difference is that a deadline cannot be violated.
 - Weight parameters. A weight is a cost / reward for executing an activity or a priority of an activity.
 - Setup time. The setups are operations performed in order to prepare resources for executing some activities. A setup is measured by the time needed to perform it (called setup time), and additionally by the cost of such an operation. The two types of setup times are: sequence-independent and sequence-dependent. A sequence-dependent setup time varies depending on the sequence of activities performed on a given resource unit, a sequence-independent setup time is the same for all possible sequences of activities.
 - Preemptibility. Each of the project activities may be either preemptable or non-preemptable.

3. VARIANTS OF THE PROJECT SCHEDULING PROBLEMS

The main variants of PSP are: Resource-constrained Project Scheduling Problem (RCPSPP), Multi-Mode Resource-Constrained Project Scheduling Problem (MRCPSPP), Construction Project Scheduling Problem (construction PSP), Project Scheduling Problem for Software Development (software PSP), Payment Scheduling Problem, and Time/cost trade-off Problem (TCTP).

1. Resource-constrained Project Scheduling Problem or RCPSPP [22]. RCPSPP is a problem of allocation of schedules that consists of scheduling a series of activities at a time determined with

limited resources, the activities are contemplated as renewable and nonrenewable resources that they are assigned in the course of the project [5]. RCPSP belongs to the class of the NP-hard problems [5]. The mathematical model of RCPSP [23] is formed by the equations 1-4: The equation 1 contains the objective function that minimizes the project duration (makespan) determined by the completion time (start time or S , since $d_{n+1} = 0$) of the dummy end activity $n + 1$. Constraints (equation 2) enforce the precedence constraints between activities. Constraints (equation 3) ensure for each resource k and each time period t , that the resource demand of the set of activities in process $P(t)$ at time t does not exceed the capacity, T being an upper bound on the project's makespan. Where: S is the start time, n is the number of activities in the project, $n + 1$ is the unique dummy end activity, d_i is the duration of activity j , i and j are A activities, k are the resources, K is the number of resources types, t is a time period, r are the units of resource, R is the resource demand of a set of activities, $P(t)$ is a process at time t , T is an upper bound on the project duration.

$$\text{Min } S_{n+1} \quad (1)$$

$$\text{subject to: } S_j - S_i \geq d_i, \quad (i, j) \in A \quad (2)$$

$$\sum_{P(t)} r_{ik} \leq R_k, \quad t = 1, \dots, T; \quad k = 1, \dots, K \quad (3)$$

$$S_i \geq 0 \quad \forall i \quad (4)$$

2. Multi-Mode Resource-Constrained Project Scheduling Problem or MRCPSP [24]. The MRCPSP involves the selection of a time/resource combination for each activity such that the total project makespan is minimized. The MRCPSP is an extension of the well-known single-mode RCPSP. The problem of finding a feasible solution for the MRCPSP is NP-Complete problem [25]. The mathematical model of MRCPSP [24] is formed by the equations 5-10: Constraints (equation 6) ensure that each non-preemptable activity is performed exactly once in exactly one mode. Precedence constraints are guaranteed by (equation 7). Constraints (equation 8 and 9) ensure that the renewable and nonrenewable resource limits are not exceeded, respectively. Finally, the equation 10 is the binary status of the decision variables. Where: $R_k^P (R_k^v)$ is the number of available units of the k th (l th) renewable (nonrenewable) resource, r_{jmk}^P is the number of units of the k th renewable resource ($k = 1, \dots, R$) required by activity A_j executed in mode $m \in M_j$, r is the number of units of the l th nonrenewable resource ($l = 1, \dots, N$) required by activity A_j executed in mode $m \in M_j$, d_{jm} is the duration of activity A_j executed in mode $m \in M_j$, $x_{jmt} = 1$ if activity A_j executed in mode $m \in M_j$ is completed at the end of time period t , EF_j , LF_j are calculated assuming that the shortest duration mode is assigned to each activity, and the planning horizon H is calculated for the modes with the longest durations.

$$\text{Minimize } \sum_{t=EF_{n+1}}^{LF_{n+1}} t x_{n+1,m,t} \quad (5)$$

$$\text{subject to } \sum_{m=1}^{|M_j|} \sum_{t=EF_j}^{LF_j} x_{jmt} = 1 \text{ for } j = 0, \dots, n+1 \quad (6)$$

$$\sum_{m=1}^{|M_j|} \sum_{t=EF_j}^{LF_j} tx_{imt} \leq \sum_{m=1}^{|M_j|} \sum_{t=EF_j}^{LF_j} tx_{jmt} - d_{jm} \text{ for all } (A_i, A_j) \in P \quad (7)$$

$$\sum_{j=1}^n \sum_{m=1}^{|M_j|} \sum_{q=\max\{t, EF_j\}}^{\min\{t+d_{jm}-1, LF_j\}} r_{jmk}^p x_{jm q} \leq R_k^p \text{ for } k=1, \dots, R; t=1, \dots, H \quad (8)$$

$$\sum_{j=1}^n \sum_{m=1}^{|M_j|} \sum_{t=EF_j}^{LF_j} r_{j ml}^v \leq R_k^p \text{ for } l=1, \dots, N \quad (9)$$

$$x_{jmt} \in \{0,1\} \text{ for } i=0, \dots, n+1; m \in M_j; t = EF_j, \dots, LF_j \quad (10)$$

3. Construction Project Scheduling Problem, construction PSP or CPSP [26]. The construction PSP consists of to select resources (workforce, machines) and their works to construction projects (for example scheduling of construction projects of highway construction scheduling), take into account particular conditions (technological and organizational methods, constraints resource availability to develop, and the optimal project schedule) where optimal means the accepted evaluation criteria (project duration, cost, quality) are met. The mathematical model of construction PSP [27] is formed by the equations 11-16: In the equation (11) is the objective function which minimizes the project completion time defined by minimizing the finish time of the unique dummy end activity $N + 1$. Constraints (12) ensure that no activity can be started until all its predecessors have been completed, and the dummy start activity 1 is assigned a value of 0. d_j in constrained conditions 12, it is any selected from the set of satisfying the conditions (13) above. During any time interval $(t - 1; t]$, there-into the constraints (14) represent that the resource utilization does not exceed the resource availability levels for any of the deterministic resource types, constraint set (15) and (16) denote that chance-constrained of resource utilization do not exceed the resource availability level for any of the fuzzy and random resource types respectively. Where: f is the finish time, $N+1$ is the unique dummy end activity, N is the number of activities in the project, d_j is the duration of activity j ($j \in UD$), f_j is the completion time (duration) of activity j , H is the set of pairs of activities indicating finish-start precedent relation, β_j is the probability or possibility for pre-given, i is an activity, k are the resources, K is the number of resources types, t is a time period, r are the units of resource, r_{ikt} is the amount of resource type k that is required by activity i at time t , R is the resource demand of a set of activities, R_{kt} is the total availability of deterministic resource type k at time t , a_{kt} is the total availability of fuzzy resource type k at time t , $b_{kt}(\xi)$ is the total availability of random resource type k at time t , ch , Pos is the possibility of fuzzy events, Pr is the probability of random events, SR is the set of total available resource requirements at the disposal of the project management is a random variable(process) with pre-given density function throughout the scheduling horizon, FR is the set of total available resource requirements at the disposal of the project management is a fuzzy number(be related to time) with pre-given membership function throughout the scheduling horizon, DR is the set of the total available resource requirements at the disposal of the project management is pre-given and fixed (for example R_{kt}) at time t , UD is the set of activity with random duration UD_1 or fuzzy duration UD_2 such that $UD = UD_1 + UD_2$.

$$\text{Minimize } f_{N+1} \quad (11)$$

$$f_j - d_j \geq f_i \text{ and } f_1 = 0, \forall (i, j) \in H \quad (12)$$

$$Ch\{d'_j \leq d_j \leq d_j^*\} \geq \beta_j, j \in UD \quad (13)$$

$$\sum_{i \in S_i} r_{ikt} \leq R_{kt}, k \in DR \quad (14)$$

$$Pos\left\{\sum_{i \in S_i} r_{ikt} \leq a_{kt}\right\} \geq \alpha_{kt}, k \in FR \quad (15)$$

$$Pr\left\{\sum_{i \in S_i} r_{ikt} \leq b_{kt}(\xi)\right\} \geq p_{kt}, k \in SR, \text{ for } t = 1, 2, \dots, f_N \quad (16)$$

4. Project Scheduling Problem for Software Development, software PSP or PSPSD [28]. The software PSP consists of human resource allocation to the various tasks in a software development project in accordance with their skills to produce higher quality software, while keeping effort expenditure and schedule time to a minimum [29]. The mathematical model of software PSP [29] is formed by the equations 17-21: the equation (17) is the objective function of the software PSP, the equation (18) contains aspects related to project development cost, the equation (19) contains aspects related to project development duration, the equation (20) contains aspects related to task precedence relations of two consecutive task of the software development project, and the equation (21) contains aspects related to personnel availability. Where: Pc is the cost per unit time of the j^{th} personnel, m is the total number of personnel allotted to the project development, f_1 is the total development cost, P_{Tj} is the processing time of the j^{th} task in the project, Q_{Tj} is the queuing time of the j^{th} task in the project for resource availability, f_2 is the actual duration of the project, S_{Tj} is the starting time of the j^{th} task, F_{Tj} is the finish time of the j^{th} task.

$$\text{minimize } f(x) = f_1 + f_2 \quad (17)$$

$$f_1 = \sum_{j=1}^m Pc_j \quad (18)$$

$$f_2 = \sum_{j=1}^m (Q_{Tj} + P_{Tj}) \quad (19)$$

$$F_{T(j-1)} < S_{Tj} \quad \forall j \quad (20)$$

$$P_j(t) \Rightarrow \tau_i(t) \neq P_j(t) \Rightarrow \tau_k(t) \quad \forall i, j, k \quad (21)$$

5. Payment Scheduling Problem [30,31]. The Payment Scheduling Problem PSP consists of are several issues associated with analyzing trade-off in contract prices, profit margins, project deadlines and other payment parameters by developing models and solution methods for simultaneously determining the among, timing and location of payments in projects. The mathematical model of payment scheduling problem [32] is formed by the equations 22-24: The equation (22) is the objective function represents cash outflows, cash inflow and capital cost, where each component is discounted back to the beginning of the project, the equation (23) are the activity precedent constraints, the equation (24) are the capital constraints for each period of the project. It assumes that capital is a renewable resource, where the initial capital available c_0 is augmented or reduced by cash flows that occur throughout the project. Where: c_0 is the total capital available at the beginning of the project in period 0, l_k is the capital investment required by activity k ($k = 1, 2, \dots, m$), $F_{i(k)}$ is the cash outflow at the beginning of activity k at node i , where each activity is defined by nodes i and j , $F_{j(k)}$ is the cash

inflow received upon completion of activity k at node j , d_k is the duration of activity k , where k may not be preempted, $T_{i(k)}$ is the time at which node i of activity k is scheduled to occur, $Z_{a(t)}$ is the set of activities (a) that are scheduled to be active in period t , $Z_{p(t)}$ is the set of activities (p) completed prior to period t , α is an opportunity cost of capital.

$$\text{Maximize } \sum_{k=1}^m F_{i(k)} \exp(-\alpha T_{i(k)}) + F_{j(k)} \exp(-\alpha T_{j(k)}) - [I_k \exp(-\alpha d_k) - I_k] \exp(-\alpha T_{i(k)}) \quad (22)$$

$$T_{j(k)} - T_{i(k)} \geq d_k \cdot k = 1, 2, \dots, m \quad (23)$$

$$\sum_{k \in Z_{a(t)}} I_k \leq c_0 + \sum_{k \in Z_{p(t)}} (F_{j(k)} + F_{i(k)}) \quad (24)$$

6. Time/cost trade-off Problem or TCTP [33,34]. The TCTP consists [35] consists of an algorithm for efficiently shortening the duration of a project when it exceeds the predetermined limit. Suppose that we are given a project network which represents with a set of activities (and their precedence relationships) to be performed (individual activity in one of several ways), each activity with its unique time and cost requirements. Suppose that we are given a project network which represents with a set of activities (and their precedence relationships) to be performed (individual activity in one of several ways), each activity with its unique time and cost requirements. Different decisions as to how the various activities are performed lead to different time-cost realizations for the overall network. The objective of the time-cost tradeoff problem for a project network is to identify the sets of decisions that result in desirable time-cost realizations. The mathematical model of $P-C|T$ [36] is formed by the equations 25-28: The equation (25) reflects the cost minimization objective, equation (26) ensures that exactly one alternative is chosen for each activity, equation (27) maintains the precedence relationships among the activities; and equation (28) guarantees that the project will complete by its due-date. Where: i are the activities, n are the nodes, j is the alternative, $a(i)$ is an alternative of activity, c_{ij} is the cost, x_{ij} is a 0-1 variable which is 1 if alternative j is selected for executing activity i and 0 otherwise, t_{ij} is the time, s_i ($s_i \geq 0$) is the start time for activity i , $S(i)$ is the set of the immediate successors of i , $n+1$ is the finish node, d is the due date.

$$\min \sum_{1 \leq i \leq n} \sum_{1 \leq j \leq a(i)} c_{ij} x_{ij} \quad (25)$$

$$\sum_{1 \leq j \leq a(i)} x_{ij} = 1 \text{ for all } i = 1, \dots, n \quad (26)$$

$$\sum_{1 \leq j \leq a(i)} t_{ij} x_{ij} + s_i \leq s_k \text{ for all } k \in S(i) \quad (27)$$

$$s_{n+1} \leq d \quad (28)$$

4. ALGORITHMS TO SOLVE THE PROJECT SCHEDULING PROBLEMS

In the specialized literature there exist various approaches for carrying out the solution of the Project Scheduling Problems by means algorithms. Table 1 shows diverse algorithms to solve the Project Scheduling Problems and their variants. Where: 1. The Tabu search algorithm applied to the RCPSP

and their variants were used by [37,38,39]. 2. The Simulated Annealing algorithm applied to the RCPSP and their variants were used by [40,41,42]. 3. The Genetic Algorithm applied to the RCPSP and their variants were used by [43, 44, 45, 46, 47, 48, 49, 50]. 4. Others algorithms (branch-and-bound) applied to the RCPSP and their variants were used by [51, 52, 53]. 5. The Tabu Search algorithm applied to the MRCPSP and their variants were used by [54, 55]. 6. The Simulated Annealing algorithm applied to the MRCPSP and their variants were used by [41, 57, 58, 59, 61,62]. 7. The Genetic Algorithm applied to the MRCPSP and their variants were used by [61, 63, 64, 65, 66, 67, 68]. 8. Others algorithms (branch-and-bound) applied to the MRCPSP and their variants were used by [51]. The Particle Swarm Optimization by [71, 72]. 9. The Simulated Annealing algorithm applied to the construction PSP and their variants were used by [73, 74]. 10. The Genetic Algorithm applied to the construction PSP and their variants were used by [74, 76]. 11. The Genetic Algorithm applied to the software PSP and their variants were used by [77, 78]. 12. Others algorithms (Particle Swarm Optimization) applied to the software PSP and their variants were used by [29]. 13. The Tabu search algorithm applied to the Payment Scheduling Problem and their variants was used by [79]. 14. The Simulated Annealing algorithm applied to the Payment Scheduling Problem and their variants was used by [61, 79]. 15. The Genetic Algorithm applied to the Payment Scheduling Problem and their variants was used by [80]. 16. Others algorithms (Hybrid Simulated Annealing and Genetic Algorithm) applied to the Payment Scheduling Problem and their variants were used by [80]. 17. The Tabu search algorithm applied to the TCTP and their variants were used by [81]. 18. Genetic Algorithm applied to the TCTP and their variants were used by [82, 83, 84]. 19. Others algorithms (branch-and-bound) applied to the TCTP and their variants were used by [85]. The hybrid scatter-search [86].

The RCPSP variants: Single-Mode Resource-Constrained Project Scheduling Problem [79], high school timetabling [87, 88], University Course Timetabling [87, 88], the audit-scheduling problem [89]. The MRCPSP variants: Multi-mode resource-constrained project scheduling problem with discounted cash flows or MRCPSPDCF [90, 91], multi-skill project scheduling problem or MSPSP [92, 93]. The construction PSP variants: larger-scale construction project scheduling problem with serious resource conflict [73], and Multi-Site Construction Problem Scheduling [95]. The software PSP variants: Fuzzy project scheduling system for software development [96], Time depend software Project Scheduling Problem [56], A project scheduling problem with labour constraints and time-dependent activities requirements [60], Project Scheduling Problem for Software Development with Random Fuzzy Activity Duration Times [28], and Fuzzy resource-constrained project scheduling problem for software development [75]. The Payment Scheduling Problem variants: Contractor's payment scheduling problem [69], client's payment scheduling problem [80], Joint payment scheduling Problem [94], max-NPV project scheduling problem [70], multi-mode project payment scheduling problem or MPPSP [62]. The TCTP variants: deadline problem [106], budget problem [36], time-cost tradeoff problem in fuzzy environment [15].

Table 1 algorithms to solve project scheduling problems and their variants.

Algorithms / methods	RCPS P and variants	MRCPS P and variants	Construction PS P and variants	Software PS P and variants	Payment Scheduling Problem	TCTP and variants
Tabu Search	1	5	-	-	13	17
Simulated Annealing	2	6	9	-	14	-
Genetic Algorithm	3	7	10	11	15	18
Others	4	8	-	12	16	19

5. CONCLUSIONS

This paper was a guide to understand Project Scheduling Problem (PSP) and their variants, contains a summary of the related works of the PSP, it provides a comprehensive view of the state of research and practice in the domain of PSP, diverse options appeared to solve PSP with costs, resources, time, and so on. We find in the related works that all the authors mention the term Project Scheduling Problem as the same that the term Resource-constrained Project Scheduling Problem, but we think necessary to make a difference with a new definition (section 2 of this paper) of the Project Scheduling Problem as a generic name given to a whole class of problems, including RCPS P and other problems. Finally, we conclude that, it's necessary to create new (others) mathematical models of real problems related to PSP, and more algorithms, techniques and applications that could solve the real problems, for example, to solve the scheduling of innovation and technology projects and scheduling of government projects, and others. The future works of this research are a focus to try to solve these problems.

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